

W Coxeter group $W = \langle s \in S \mid (st)^{m_{st}} = 1, s^2 = 1 \rangle$

\mathbb{K} field

$W \subset h / \mathbb{K}$ "reflection representation", i.e. $\{\alpha_s\} \subseteq h^*, \{\alpha_s^\vee\} \subseteq h$ $m_{st} \in \{2, 3, \dots, \infty\}$

$s(v) = v - \alpha_s^\vee(v)\alpha_s$ defines a representation of W

TECHNICAL ASSUMPTION reflection faithful i.e.

$$T = \bigcup_w S_w^{-1} \text{ reflections}$$

h^y has codim 1 $\Leftrightarrow y \in T$. (& good notion of positive roots)

CLASSICAL EXAMPLE $W = S_m \subset h^m$, char $\mathbb{K} \neq 2$

$$t = (i j) \quad \alpha_t = e_i - e_j = \alpha_{s_i} + \dots + \alpha_{s_{j-1}}$$

REFLECTION FAITHFUL $\Rightarrow t \in T \iff \alpha_t \in h^*, \alpha_t^\vee \in h$
 s.t. $t(v) = v -$

Out of (W, h) we construct the moment graph.

Vertices $x \in W$

Edges $x \xrightarrow{\alpha_t} tx \quad \forall x \in W \quad t \in T \text{ s.t. } tx \geq x$

E.g. $W = S_3 = \langle s, t \rangle$

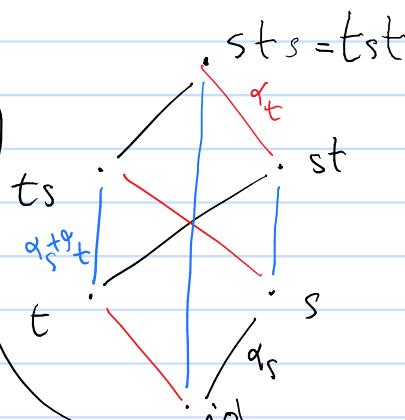
$$R = \text{Sym}_W(h^*) \cdot \deg(h^*) = 2$$

Only a sheaf M on the moment graph
 is the data of

- M_x graded R -module $\forall x \in W$

- $M_{x \rightarrow tx} \cong M_x$ "edge s.t. $\alpha_t M_{x \rightarrow tx} = 0$

- morphism $M_x \xrightarrow{p_{x,tx}} M_{x \rightarrow tx} \leftarrow M_{tx}$



EXAMPLE Constant sheaf A : $A_x = R$, $A_{x-tx} = R/(g_t)$

$$\text{Global sections } \Gamma(A) = \left\{ (m_x)_{x \in W}^M \mid \begin{array}{l} M_x \\ \downarrow \rho \\ M_{tx} \end{array} \right. \quad \left. \begin{array}{l} M_{tx} \\ \downarrow \rho \\ p_{tx}(m_x) = p_{tx}(m_{tx}) \end{array} \right\}$$

$\mathbb{Z} := \Gamma(A)$ is a ring

U1

\mathbb{Z} bounded sections (i.e. $\exists k$ s.t. $\deg(a_x) \leq k$)

Thm If W is a Weyl group (of Lie-Moody group G) then

$$\mathbb{Z} \cong H^*(G/B, \mathbb{K}).$$

For general Coxeter groups \mathbb{Z} is free with basis indexed $\{\xi_x\}_{x \in W}$ and $\deg(\xi_x) = 2\ell(x)$

$\mathbb{Z} \cong$ constant-hyper dual nil Hecke ring.

Moment graphs provide a combinatorial setting for studying
intersection cohomology sheaves (Braden-MacPherson '01)
or Parity sheaves (Fiebig-Williams '11)

Braden-MacPherson algorithm.

(Construction of the canonical sheaves $B(x)$)

Start with $(B(x))_x = R$, $(B(x))_z = 0$ if $z \notin x$.

Ann $B(x)_z$ defined for all $z > y$

$$E = \{y \rightarrow t_y\} \quad B(x)_{y \rightarrow t_y} = B(x)_{t_y} / (g_t) \quad \text{for all arrows}$$

$$\mathcal{P}(\mathcal{B}_x, \geq_y) \xrightarrow{\cong} \bigoplus_{e \in E} \mathcal{B}_e$$

\in

$(\mathcal{B}_x)_y := \text{proj. cover of } \text{Im } q$

(in other words): $(\mathcal{B}_x)_y$ is free and induces isomorph.

$$(\mathcal{B}_x)_y \otimes_R \mathbb{K} \xrightarrow{\sim} (\text{Im } q) \otimes_R \mathbb{K}$$

Thm (Elliott-Williamson) $\text{gr rk } (\mathcal{B}_x)_y = P_{x,y}(q)$

determine characters
of m . modules of
reductive Lie algebras in
the category \mathcal{O}

\leftarrow LC polynomial

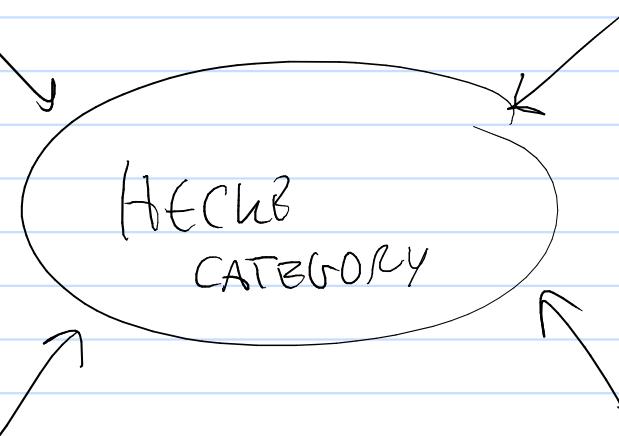
Thm (Fiebig-Williamson, '11)
 \mathfrak{W} Weyl group, char \mathbb{K}^P

$\text{gr rk } \mathcal{B}(x)_y = \text{pk } P$ polynomials

↓
determines character of tilting
modules of reductive alg. group
in char P

BMP shuffles
on the moment
graphs

Songel Lascouxes



Parity shuffles
(or IC shuffles in char 0)

Elliott-Williamson
diagrammatic category.

Soergel bimodules heronian monoidal additive category
graded R -bimodule generated by $R \otimes_{R^S} R$ and shifts

Indecomposable Soergel bimodule are B_x , $x \in W$

$\Gamma(B_x)$ is the indecomposable Soergel bimodule B_x

Notice that this gives a natural structure of \mathbb{Z} -module on B_x

which $\mathbb{S}\text{Bim}$ do not see $R \otimes R \rightarrow \mathbb{Z}$ is not surjective
in general for infinite Coxeter group.

This becomes meaningful when we look at Soergel modules

B_x indecomp. Soergel bimodule

$$\overline{B}_x := B_x \otimes_R \mathbb{R}$$

Thm (P.) \overline{B}_x is not indecomposable as a R -module
but it is indecomposable as a \mathbb{Z} -module

$$\text{Ex: } w = \text{startst in } \widehat{A}_2 = \begin{matrix} & \nearrow \\ \widehat{\mathbb{Z}} & \end{matrix}$$

Thm (P.) $\widehat{\mathbb{Z}}$ is the center of the Hecke category.

COMBINATORICS & THE COEFFICIENT OF q (Assume \mathbb{R})

$x \leq y \in W$ $[x, y] = \{z \in W \mid x \leq z \leq y\}$ Bruhat interval

BMP algorithm \Rightarrow Restriction of the moment graph to $[x, y]$

\rightsquigarrow determines KL polynomial $P_{x,y}(q)$

Conj (Combinatorial invariance, Lusztig, Lycen '80s)

The KL polynomial $P_{x,y}(q)$ depends only on the graph type of $[x, y]$ (i.e. no local needed)

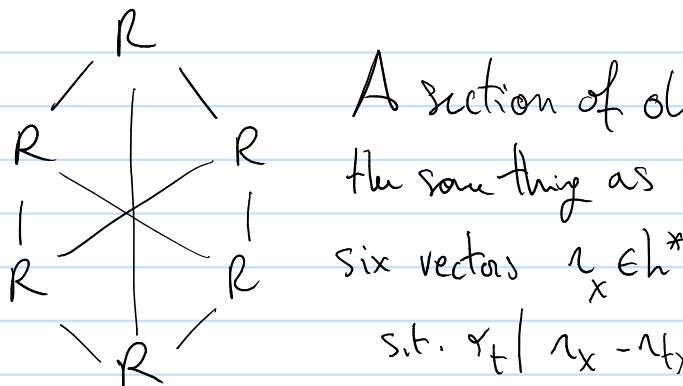
Still open! Some partial results

Brenti - Caselli - Marietti '08 $P_{e,x}(q)$ is a combinatorial invariant.

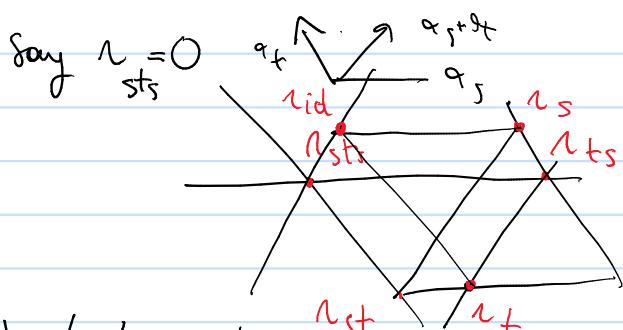
Conj holds if $\ell(x) - \ell(y) \leq 4$
 $(\leq 8 \text{ in type A})$

Still unknown even for the coefficient of q

$q_{x,y} := \text{coeff. of } q \text{ in } P_{x,y}(q) = 1 + q_{x,y} q + \text{"high terms"}$



A section of object 2 is
the same thing as choosing
six vectors $\alpha_x \in \mathbb{C}^*$
s.t. $\alpha_t \mid \alpha_x - \alpha_{t x}$



Want to make 2 points here.

- 1) n_{st}, n_{ts} determine the section. 2) the 3 lines always intersect in a point because of Pappus theorem.

$$q_{x,y} = \text{codim}\left(R_2 \xrightarrow{\pi} (\mathcal{B}^{\delta_x})_2\right)$$

$\leftarrow \text{degree 2}$

$$= \dim \Gamma(\mathcal{B}(x), >y)_2 - \dim \left\{ \begin{array}{l} \text{π-extendable} \\ \text{sections} \end{array} \right\} =: V_{x,y}$$

A way to see $V_{x,y}$: if we fix the length of the top edges, when this give rise to a solution?

In the example above $P_{sts,x}(q) = 1$ since all sections extend

$$\begin{aligned} c_{x,y} &= \# \text{ coatons in } [x,y] \\ &= \# \{ x \geq z \geq y \mid l(z) = l(x) - 1 \} \\ &= \dim \Gamma(\mathcal{B}_x, >x)^2 \end{aligned}$$

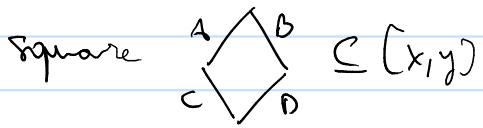
$$d_{x,y} = \dim V_{x,y}$$

$$\text{Then (Dyer '97)} \quad q_{x,y} = c_{x,y} - d_{x,y} \quad \begin{pmatrix} \text{BFS FOR MOMBUT} \\ \text{GRAPHS WORK!} \\ \text{INTRODUCTORY!} \end{pmatrix}$$

$E_{x,y}$ set of edges of $[x,y]$

\cup
 F

let $\tilde{F} \supseteq F$ minimal such that whenever we have a

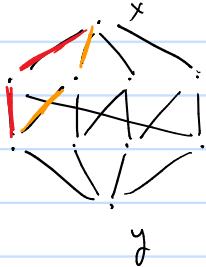


$$A, B \in \tilde{F} \Rightarrow C, D \in \tilde{F}$$

We say \tilde{F} is square-generating if $\tilde{F} = E_{x,y}$

$g_{x,y}$:= minimal size of a generating set.

EXAMPLES



Obs $g_{x,y} \geq d_{x,y}$ because fixing the length of the edges in F determines a section.

Obs. $g_{x,y} \leq c_{x,y}$. top edges always generate

$g_{x,y} \leq l(y) - l(x)$: every maximal chain generates

(follows from shellability But what interval)

Thm(P.) In type A $g_{x,y} = d_{x,y}$. (part of a longer joint project w/ Williamson)

Cor $q_{x,y} = c_{x,y} - g_{x,y}$ is combinatorial invariant in type A

Thm Proof uses crucially generalized lifting property
(Tsimkovich-Williams '15)

$x \leq y \exists$ reflection t s.t. $x < tx < y$ and $x < ty < y$ and
such that $R_{x,y}(q) = (q-1)R_{x,ty} + qR_{tx,ty}$.